CS 70 Discrete Mathematics and Probability Theory
Spring 2017 Rao

HW 5

1 Sundry

Before you start your homework, write down your team. Who else did you work with on this homework? List names and email addresses. (In case of homework party, you can also just describe the group.) How did you work on this homework? Working in groups of 3-5 will earn credit for your "Sundry" grade.

Please copy the following statement and sign next to it:

I certify that all solutions are entirely in my words and that I have not looked at another student's solutions. I have credited all external sources in this write up.

2 Count and Prove

(a) Over 1000 students walked out of class and marched to protest the war. To count the exact number of students protesting, the chief organizer lined the students up in columns of different length. If the students are arranged in columns of 3, 5, and 7, then 2, 3, and 4 people are left out, respectively. What is the minimum number of students present? Solve it with Chinese Remainder Theorem.

(b) Prove that for $n \geq 1$, if $935 = 5 \times 11 \times 17$ divides $n^{80} - 1$, then 5, 11, and 17 do not divide n.

3 RSA Lite

Woody misunderstood how to use RSA. So he selected prime $P = 101$ and encryption exponent $e = 67$, and encrypted his message m to get $35 = m^e \bmod P$. Unfortunately he forgot his original message m and only stored the encrypted value 35. But Carla thinks she can figure out how to recover m from $35 = m^e \bmod P$, with knowledge only of P and e. Is she right? Can you help her figure out the message m? Show all your work.

4 RSA with Three Primes

Show how you can modify the RSA encryption method to work with three primes instead of two primes (i.e. $N = pqr$ where p, q, r are all prime), and prove the scheme you come up with works in the sense that $D(E(x)) \equiv x \bmod N$.

5 Squared RSA

(a) Prove the identity $a^{p(p-1)} \equiv 1 \pmod{p^2}$, where a is relatively prime to p and p is prime.

(b) Now consider the RSA scheme: the public key is $(N = p^2q^2, e)$ for primes p and q, with e relatively prime to $p(p-1)q(q-1)$. The private key is $d = e^{-1} \pmod{p(p-1)q(q-1)}$. Prove that the scheme is correct, i.e. $x^{ed} \equiv x \pmod{N}$.

(c) Continuing the previous part, prove that the scheme is unbreakable, i.e. your scheme is at least as difficult as ordinary RSA.

6 Breaking RSA

(a) Eve is not convinced she needs to factor $N = pq$ in order to break RSA. She argues: "All I need to know is $(p-1)(q-1)$... then I can find d as the inverse of $e \bmod (p-1)(q-1)$. This should be easier than factoring N." Prove Eve wrong, by showing that if she knows $(p-1)(q-1)$, she can easily factor N (thus showing finding $(p-1)(q-1)$ is at least as hard as factoring N). Assume Eve has a friend Wolfram, who can easily return the roots of polynomials over $\mathbb{R}$ (this is, in fact, easy).

(b) When working with RSA, it is not uncommon to use $e = 3$ in the public key. Suppose that Alice has sent Bob, Carol, and Dorothy the same message indicating the time she is having her birthday party. Eve, who is not invited, wants to decrypt the message and show up to the party. Bob, Carol, and Dorothy have public keys $(N_1, e_1), (N_2, e_2), (N_3, e_3)$ respectively, where $e_1 = e_2 = e_3 = 3$. Furthermore assume that N_1, N_2, N_3 are all different. Alice has chosen a number $0 \leq x < \min\{N_1, N_2, N_3\}$ which indicates the time her party starts and has encoded it via the three public keys and sent it to her three friends. Eve has been able to obtain the three encoded messages. Prove that Eve can figure out x. First solve the problem when two of N_1, N_2, N_3 have a common factor. Then solve it when no two of them have a common factor. Again, assume Eve is friends with Wolfram as above.

Hint: The concept behind this problem is the Chinese Remainder Theorem: Suppose $n_1, ..., n_k$ are positive integers, that are pairwise co-prime. Then, for any given sequence of integers

$a_1, ..., a_k$, there exists an integer x solving the following system of simultaneous congruences:

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ x &\equiv a_2 \pmod{n_2} \\ &\vdots \\ x &\equiv a_k \pmod{n_k} \end{aligned}$$

Furthermore, all solutions x of the system are congruent modulo the product, $N = n_1 \cdots n_k$. Hence: $x \equiv y \pmod{n_i}$ for $1 \leq i \leq k \Leftrightarrow x \equiv y \pmod{N}$.

7 Polynomials in Fields

Define the sequence of polynomials by $P_0(x) = x + 12$, $P_1(x) = x^2 - 5x + 5$ and $P_n(x) = xP_{n-2}(x) - P_{n-1}(x)$.

(For instance, $P_2(x) = 17x - 5$ and $P_3(x) = x^3 - 5x^2 - 12x + 5$.)

(a) Show that $P_n(7) \equiv 0 \pmod{19}$ for every $n \in \mathbb{N}$.

(b) Show that, for every prime q, if $P_{2017}(x) \not\equiv 0 \pmod{q}$, then $P_{2017}(x)$ has at most 2017 roots modulo q.